



# Solving the Minimum Spanning Tree Problem Under Interval-Valued Fermatean Neutrosophic Domain

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**Abstract:** In classical graph theory, the minimal spanning tree (MST) is a subgraph that lacks cycles and efficiently connects every vertex by utilizing edges with the minimum weights. The computation of a minimum spanning tree for a graph has been a pervasive problem over time. However, in practical scenarios, uncertainty often arises in the form of fuzzy edge weights, leading to the emergence of the Fuzzy Minimum Spanning Tree (FMST). This specialized approach is adept at managing the inherent uncertainty present in edge weights within a fuzzy graph, a situation commonly encountered in real-world applications. This study introduces the initial optimization approach for the Minimum Spanning Tree Problem within the context of interval-valued fermatean neutrosophic domain. The proposed solution involves the adaptation of the Dhouib-Matrix-MSTP (DM-MSTP) method, an innovative technique designed for optimal resolution. The DM-MSTP method operates by employing a column-row navigation strategy through the adjacency matrix. To the best of our knowledge, instances of this specific problem have not been addressed previously. To address this gap, a case study is generated, providing a comprehensive application of the novel DM-MSTP method with detailed insights into its functionality and efficacy.

**Keywords:** Minimum Spanning Tree Problem; Fermatean Neutrosophic Domain; Dhouib-Matrix-MSTP; Artificial Intelligence; Operations Research; Combinatorial Optimization

## 1. Introduction

The Minimum Spanning Tree (MST) is a fundamental concept in graph theory, serving as a crucial optimization problem with widespread applications in various fields. At its core, the MST seeks to identify the most efficient way to connect a set of vertices within a graph by utilizing edges with minimal total weights. This problem has garnered significant attention due to its relevance in operational research, communication systems, transportation networks, logistics, supply chain management, image processing, wireless telecommunication, and cluster analysis. Consider a connected graph  $G=(V,E)$  where  $V$  represents the set of vertices and  $E$  is the set of edges. In the context of this graph, a tree denoted as  $T$  is considered a spanning tree if it serves as a subgraph of  $G$  and

encompasses all the nodes present in  $G$ . Referred to as a maximal tree subgraph,  $T$  stands out among other trees within  $G$  as the most extensive, containing the maximum number of arcs.

Conventional spanning trees operate under the assumption of precise and deterministic edge weights, which may not align with the realities of many real-world situations characterized by ambiguous or imprecise information. The introduction of fuzzy logic offers a valuable means to address uncertainty, providing a more accurate model for scenarios where the nature of relationships between nodes is not fully understood or easily quantifiable. The necessity for a Fuzzy Spanning Tree (FST) arises as a response to the limitations of traditional spanning tree models in effectively handling the inherent uncertainties found in practical, real-world scenarios. Recognizing that real-world systems often grapple with imprecise, uncertain, or incomplete information, the adoption of fuzzy logic in spanning tree models emerges as a solution. This approach enhances the models' ability to capture and represent inherent uncertainties, resulting in solutions that are more robust and adaptable across various application domains. Zadeh [34] introduced the concept of fuzzy sets to address ambiguity and uncertainty by employing a degree of membership function. Subsequently, Atanassov [35] proposed intuitionistic fuzzy sets (IFS), capable of managing both membership and non-membership functions, offering enhanced flexibility in dealing with ambiguity.

Building upon these ideas, Smarandache [36] developed the concept of neutrosophic sets (NS), which includes membership, indeterminacy, and non-membership functions. NS gained considerable attention, leading to numerous studies [37]-[41]. In the domain of minimum spanning trees (MST), Agnes et al. [29] formulated a Fuzzy Clustering model, while Oscan et al. [30] addressed uncertain cost and demand parameters with a capacitated fuzzy MST approach. Mohanta et al. [31] proposed an algorithm for intuitionistic fuzzy MST, and Mandal et al. [32] explored robust MST in cancer detection using intuitionistic fuzzy graphs. For neutrosophic fuzzy graphs, Dey et al. [33] implemented an algorithm for MST.

Senapati et al. [42] first introduced Fermatean fuzzy sets as an extension of Pythagorean fuzzy sets to overcome their inherent limitations. Following this, Jeevaraj [43] further developed the concept by proposing interval-valued Fermatean fuzzy sets. Subsequently, Palani et al. [44] elaborated on the topic by presenting a decision-making methodology within the framework of Pythagorean interval-valued Fermatean fuzzy sets. Expanding on this research, Ruan et al. [45] discussed a multi-criteria decision-making approach tailored for interval-valued Fermatean neutrosophic fuzzy sets.

Despite these advancements, there has been a notable absence of matrix-based methods for MST in the literature. This gap in existing methods has motivated our work, leading us to extend the concept to the Dhouib-Matrix-MSTP (DM-MSTP) method for interval-valued fermatean fuzzy graphs.

This paper is structured as follows: Section 2 introduces the fundamental concepts of IFS, NFS, FFS, and IVFNS. Section 3 outlines the Dhouib Matrix MST method for IVFNN. Section 4 provides numerical examples, and Section 5 concludes the paper.

## 2. Preliminaries

**Definition 1.** The Fermatean fuzzy Set (FFS)  $\tilde{F}$  in the universal set  $X$  is defined by  $\tilde{F} = \{(x, \mu_{\tilde{F}}(x), \nu_{\tilde{F}}(x)) : x \in X\}$  where the membership function  $\mu_{\tilde{F}}(x) : X \rightarrow [0, 1]$  and the non-membership

function  $\nu_F(x): X \rightarrow [0, 1]$  satisfy the condition  $[\mu_F(x)]^3 + [\nu_F(x)]^3 \leq 1$  is said to be the degree of hesitation of  $x$  to  $\tilde{F}$ .

**Definition 2.** Let  $X$  be the universe of discourse. Then  $N = \{(x, T_N(x), I_N(x), F_N(x)): x \in X\}$  is defined as Neutrosophic Fuzzy Set (NFS), where the truth-membership function is represented as  $T_N(x): X \rightarrow [0,1]$  an indeterminacy-membership function  $I_N(x): X \rightarrow [0,1]$  and the falsity membership function  $F_N(x): X \rightarrow [0,1]$  which satisfies the conditions  $0 \leq T_N(x) + I_N(x) + F_N(x) \leq 3, \forall x \in X$ .

**Definition 3.** A neutrosophic fuzzy set  $\ell$  in the universe  $X$  is the form of  $\ell = \{(u, T_\ell(u), I_\ell(u), F_\ell(u)): u \in \ell\}$  represents the degree of truth, indeterminacy and falsity-membership of  $\ell$  respectively. The mapping  $T_\ell(u): \ell \rightarrow [0,1], I_\ell(u): \ell \rightarrow [0,1], F_\ell(u): \ell \rightarrow [0,1]$  and  $0 \leq T_\ell(u)^3 + I_\ell(u)^3 + F_\ell(u)^3 \leq 2$ . Here  $\ell = (T_\ell, I_\ell, F_\ell)$  is denoted as fermatean neutrosophic number (FNN).

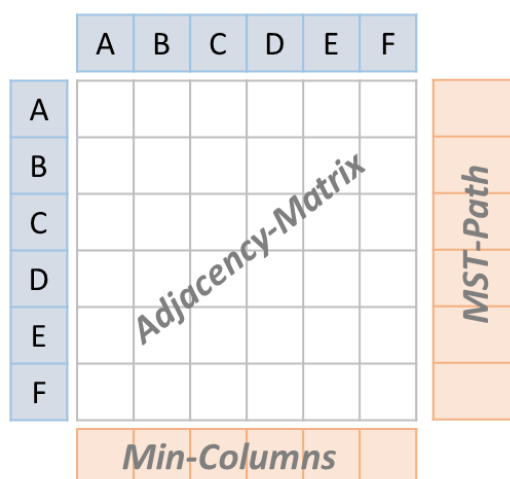
**Definition 4.** [28] An interval-valued fermatean neutrosophic set (IVFNS)  $A$  on the universe of discourse  $X$  is defined as  $A = \{(x, T_A(x), I_A(x), F_A(x)) | x \in X\}$ , where  $T_A(x) = [T_A^-(x), T_A^+(x)], I_A(x) = [I_A^-(x), I_A^+(x)], F_A(x) = [F_A^-(x), F_A^+(x)]$  represents the truth-membership degree, indeterminacy-membership degree and falsity-membership degree, respectively. Hence the mapping  $T_A(x): X \rightarrow [0,1], I_A(x): X \rightarrow [0,1], F_A(x): X \rightarrow [0,1]$  and  $0 \leq (T_A(x))^3 + (F_A(x))^3 \leq 1$  and  $0 \leq (I_A(x))^3 \leq 1$   $0 \leq (T_A(x))^3 + (F_A(x))^3 + (I_A(x))^3 \leq 2 \forall x \in X$

**Definition 5.** [28] Let  $A = \langle [T_A^L, T_A^U], [I_A^L, I_A^U], [F_A^L, F_A^U] \rangle$  then the score function  $S(x)$  is defined as:

$$S(x) = \frac{(T_A^L(x))^3 + (T_A^U(x))^3 + (I_A^L(x))^3 + (I_A^U(x))^3 + (F_A^L(x))^3 + (F_A^U(x))^3}{2}$$

### 3. The Dhouib-Matrix-MSTP method

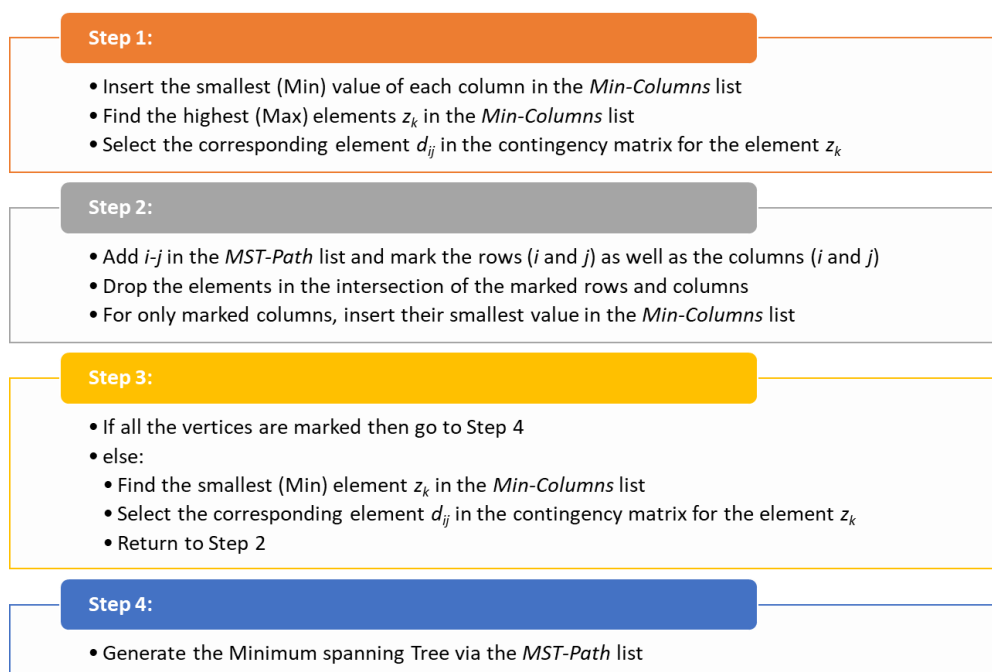
In a very recent work, the first author invented a new optimal method entitled Dhouib-Matrix-MSTP (DM-MSTP) to optimally solve the Minimum Spanning Tree Problem. The general structure of DM-MSTP is depicted in Figure 1 where two lists (the MST-Path and the Min-Columns) are added to the Adjacency-Matrix.



**Figure 1.** The general structure of DM-MSTP

Indeed, MST-Path gathers the generated components of the spanning tree, Min-Columns is used to drive the research process and Adjacency-Matrix represents the distance between all the vertices.

DM-MSTP is composed of four steps (see Figure 2) and for a more clarification a detailed example will be presented in section 4.



**Figure 2.** The general structure of DM-MSTP

DM-MSTP is designed under the general concept of Dhouib-Matrix where several optimization methods are developed such as the DM-ALL-SPP to create the shortest path between all-pairs of vertices in (Dhouib, 2024b) and the DM-SPP to solve single-pair, single-source and single-destination Shortest Path Problems (Dhouib, 2023a; Dhouib, 2023b; Dhouib, 2024c; Dhouib, 2024d). In addition, two heuristics (DM-AP1 and DM-AP2) are developed for the Assignment Problems in (Dhouib, 2022a; Dhouib, 2022b; Dhouib, 2023c; Dhouib and Sutikno, 2023) and the DM-TP1 is invented for the Transportation Problems in (Dhouib, 2021a; Dhouib, 2021b). Moreover, two other methods (DM-TSP1 and DM-TSP2) are designed for the Travelling Salesman Problems in (Dhouib, 2021c; Dhouib, 2021d; Dhouib, 2022c; Dhouib et al., 2021; Dhouib et al., 2023). Besides, three innovative metaheuristics are developed: DM4 (Dhouib, 2024e; Dhouib, 2022d; Dhouib and Pezer, 2022; Dhouib and Pezer, 2023; Dhouib, 2023d; Dhouib et al, 2024), DM3 (Dhouib, 2021e; Dhouib and Zouari, 2023a; Dhouib and Zouari, 2023b) and FtN (Dhouib, 2022e).

#### 4. Numerical examples

In this section, the DM-MSTP is simulated on an undirected graph under interval-valued Fermatean neutrosophic domains where the objective is to create the shortest minimum spanning tree. The interval-valued Fermatean neutrosophic number  $A = \langle [T_A^L, T_A^U], [I_A^L, I_A^U], [F_A^L, F_A^U] \rangle$  is converted to a crisp number using Equation 1 originally developed by (Broumi et al., 2023):

$$S(x) = \frac{(T_A^L(x))^3 + (T_A^U(x))^3 + (I_A^L(x))^3 + (I_A^U(x))^3 + (F_A^L(x))^3 + (F_A^U(x))^3}{2} \tag{1}$$

Consider the graph in Figure 3, with six vertices and nine edges.

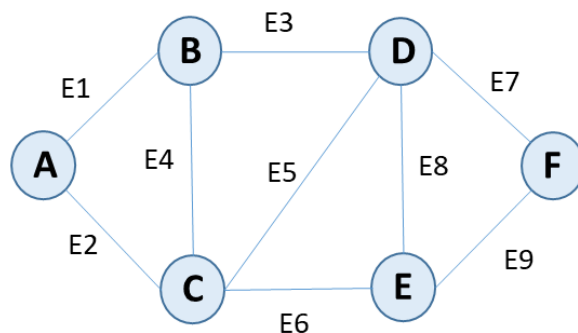


Figure 3. Undirected graph with six vertices and nine edges

The values of the nine edges are represented as interval-valued Fermatean neutrosophic numbers and are given in Table 1.

Table 1. The edge interval-valued Fermatean neutrosophic values

| Edge name | Neutrosophic length                               | Crisp length |
|-----------|---|--------------|
| E1        | $\langle [0.4,0.9], [0.5,0.6], [0.1,0.3] \rangle$ | 0.5810       |
| E2        | $\langle [0.5,0.9], [0.4,0.6], [0.1,0.5] \rangle$ | 0.6300       |
| E3        | $\langle [0.2,0.9], [0.3,0.5], [0.1,0.4] \rangle$ | 0.4770       |
| E4        | $\langle [0.2,0.6], [0.4,0.6], [0.8,0.9] \rangle$ | 0.8725       |
| E5        | $\langle [0.1,0.6], [0.5,0.8], [0.6,0.8] \rangle$ | 0.7910       |
| E6        | $\langle [0.2,0.5], [0.4,0.7], [0.6,0.8] \rangle$ | 0.6340       |
| E7        | $\langle [0.1,0.7], [0.5,0.7], [0.6,0.8] \rangle$ | 0.7700       |
| E8        | $\langle [0.3,0.9], [0.3,0.5], [0.2,0.5] \rangle$ | 0.5205       |
| E9        | $\langle [0.5,0.9], [0.2,0.4], [0.1,0.6] \rangle$ | 0.5715       |


The first step is to convert the interval-valued Fermatean neutrosophic set to a crisp one using the score function developed in Equation 1. Figure 4, illustrates the crisp matrix.

$$\begin{bmatrix}
 \infty & 0.5810 & 0.6300 & \infty & \infty & \infty \\
 0.5810 & \infty & 0.4770 & 0.8725 & \infty & \infty \\
 0.6300 & 0.4770 & \infty & 0.7910 & 0.6340 & \infty \\
 \infty & 0.8725 & 0.7910 & \infty & 0.7700 & 0.5205 \\
 \infty & \infty & 0.6340 & 0.7700 & \infty & 0.5715 \\
 \infty & \infty & \infty & 0.5205 & 0.5715 & \infty
 \end{bmatrix}$$

Figure 4. The crisp adjacency matrix

DM-MSTP starts by inserting the smallest elements of each column in Min-Columns and selecting the biggest value (0.5810) at column 1. The corresponding element (dBA) in column 1 is selected to indicate that vertices B and A are connected and 'B-A' is archived in MSTP-Path (see Figure 5).


|   | A      | B      | C      | D      | E      | F      |     |
|---|--------|--------|--------|--------|--------|--------|-----|
| A |        | 0.5810 | 0.6300 |        |        |        | B-A |
| B | 0.5810 |        | 0.4770 | 0.8725 |        |        |     |
| C | 0.6300 | 0.4770 |        | 0.7910 | 0.6340 |        |     |
| D |        | 0.8725 | 0.7910 |        | 0.7700 | 0.5205 |     |
| E |        |        | 0.6340 | 0.7700 |        | 0.5715 |     |
| F |        |        |        | 0.5205 | 0.5715 |        |     |
|   | 0.5810 | 0.4770 | 0.4770 | 0.5205 | 0.5715 | 0.5205 |     |


B-A

**Figure 5.** The vertices B and A are connected

Besides, the rows and columns (B and A) are selected (see elements with the yellow color in Figure 6) and the values of elements in the intersection are dropped. Next, Min-Columns is initiated with the smallest element of only the selected column (with the yellow color), the smallest value is selected (0.4770) and its corresponding element (dCB) is designated. Thus, vertex C is connected to vertex B and 'C-B' is archived in MSTP-Path.


|   | A      | B      | C      | D      | E      | F      |     |
|---|--------|--------|--------|--------|--------|--------|-----|
| A |        |        | 0.6300 |        |        |        | B-A |
| B |        |        | 0.4770 | 0.8725 |        |        | C-B |
| C | 0.6300 | 0.4770 |        | 0.7910 | 0.6340 |        |     |
| D |        | 0.8725 | 0.7910 |        | 0.7700 | 0.5205 |     |
| E |        |        | 0.6340 | 0.7700 |        | 0.5715 |     |
| F |        |        |        | 0.5205 | 0.5715 |        |     |
|   | 0.6300 | 0.4770 |        |        |        |        |     |


C-B

**Figure 6.** The vertices C and B are connected

Next, row and column C are selected and the value of the elements in the intersection of A, B and C are dropped (see Figure 7). Also, the smallest elements for each selected column are inserted in Min-Columns, their smallest value (0.6340) is selected and its corresponding element (dEC) is identified. Then, vertex E is connected to vertex C and 'E-C' is added to MSTP-Path.


|   | A | B      | C      | D      | E      | F      |     |
|---|---|--------|--------|--------|--------|--------|-----|
| A |   |        |        |        |        |        | B-A |
| B |   |        |        | 0.8725 |        |        | C-B |
| C |   |        |        | 0.7910 | 0.6340 |        | E-C |
| D |   | 0.8725 | 0.7910 |        | 0.7700 | 0.5205 |     |
| E |   |        | 0.6340 | 0.7700 |        | 0.5715 |     |
| F |   |        |        | 0.5205 | 0.5715 |        |     |
|   |   | 0.8725 | 0.6340 |        |        |        |     |


E-C

**Figure 7.** The vertices E and C are connected

Following, row and column E are selected and the value of the elements in the intersection are discarded (see Figure 8). Similarly, Min-Columns is initiated, the smallest value (0.5715) is selected and its corresponding element (dFE) is identified. Then, vertex F is connected to vertex E and 'F-E' is added to MSTP-Path.


|   | A | B      | C      | D      | E      | F      |     |
|---|---|--------|--------|--------|--------|--------|-----|
| A |   |        |        |        |        |        | B-A |
| B |   |        |        | 0.8725 |        |        | C-B |
| C |   |        |        | 0.7910 |        |        | E-C |
| D |   | 0.8725 | 0.7910 |        | 0.7700 | 0.5205 | F-E |
| E |   |        |        | 0.7700 |        | 0.5715 |     |
| F |   |        |        | 0.5205 | 0.5715 |        |     |
|   |   | 0.8725 | 0.7910 |        | 0.5715 |        |     |


F-E

**Figure 8.** The vertices F and E are connected

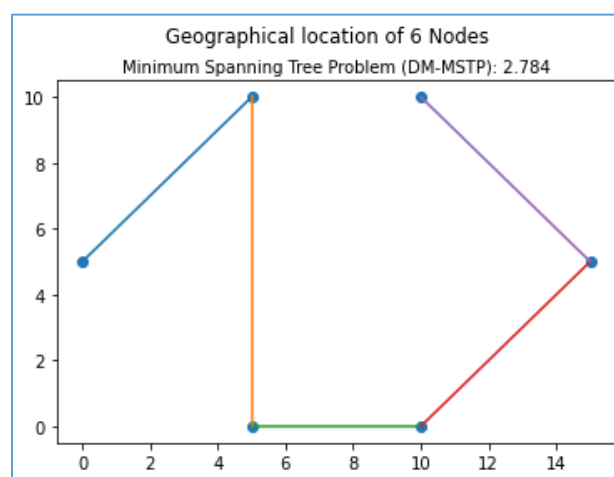
Subsequent, row and column E are selected and the value of the elements in the intersection are discarded (see Figure 9). Similarly, the Min-Columns is initiated, the smallest value (0.5205) is selected and its correspondent element is (dDF) identified. Then, vertex DE is connected to vertex F and 'D-F' is added to MSTP-Path.

|   | A | B      | C      | D      | E      | F      |     |
|---|---|--------|--------|--------|--------|--------|-----|
| A |   |        |        |        |        |        | B-A |
| B |   |        |        | 0.8725 |        |        | C-B |
| C |   |        |        | 0.7910 |        |        | E-C |
| D |   | 0.8725 | 0.7910 |        | 0.7700 | 0.5205 | F-E |
| E |   |        |        | 0.7700 |        |        | D-F |
| F |   |        |        | 0.5205 |        |        |     |
|   |   | 0.8725 | 0.7910 |        | 0.7700 | 0.5205 |     |


D-F

**Figure 9.** The vertices D and F are connected

Finally, row and column D are selected and the value of the elements in the intersection are discarded. After the above steps, the adjacency matrix is empty and the minimum spanning tree can be generated from MSTP-Path with a total cost of  $0.5810+0.4770+0.6340+0.5715+0.5205 = 2.784$ . Clearly, from Figure 10 you can see the proposed solution generated by DM-SPP.



**Figure 10.** The minimum spanning tree generated by DM-MSTP

## 5. Conclusions

This paper introduces enhancements to the Dhouib-Matrix-MSTP (DM-MSTP) method to optimize the Minimum Spanning Tree Problem within an interval-valued Fermatean neutrosophic framework. The DM-MSTP method, known for its efficiency, incorporates two key lists: the Path-Memory list, which stores spanning tree edges, and the Min-Column list, used to determine the next column to activate. Additionally, the paper includes a detailed case study illustrating the application of the method. Future research will explore the application of DM-MSTP to other variants of the Minimum Spanning Tree Problem, such as the capacitated version, and its extension to other neutrosophic domains.

**Conflicts of Interest:** The authors declare no conflict of interest.

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